

# Demystifying Blockchains: Decentralized, Secure and Fault-tolerant Storage

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In Collaboration with:

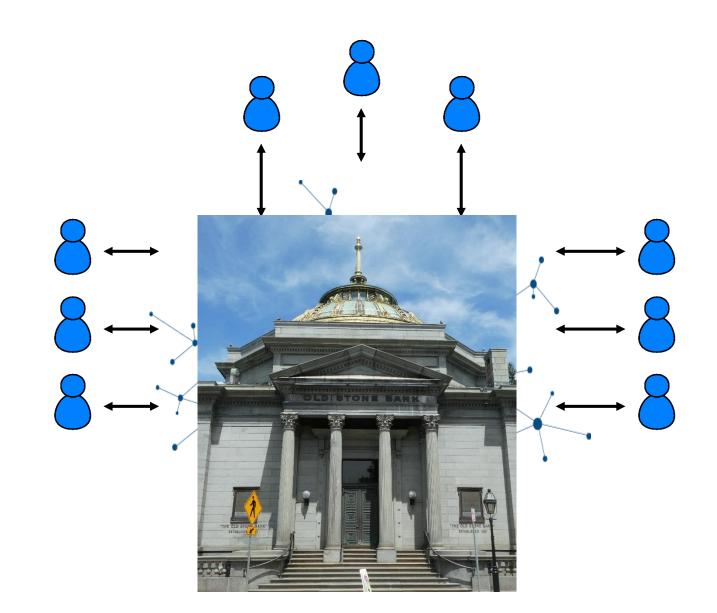
Mohammad Amiri, Sujaya Maiyya, Victor Zakhary, and **Divyakant Agrawal.** 

# Blockchains

- Many interesting (controversial?) problems in new guises.
  - Distributed Systems: Consensus, replication, etc
  - Data Management: Transactions, replication, commitment, etc
  - Security: Encryption, hashing, etc
  - Economics: Money, tokens, assests, etc



#### Bitcoin



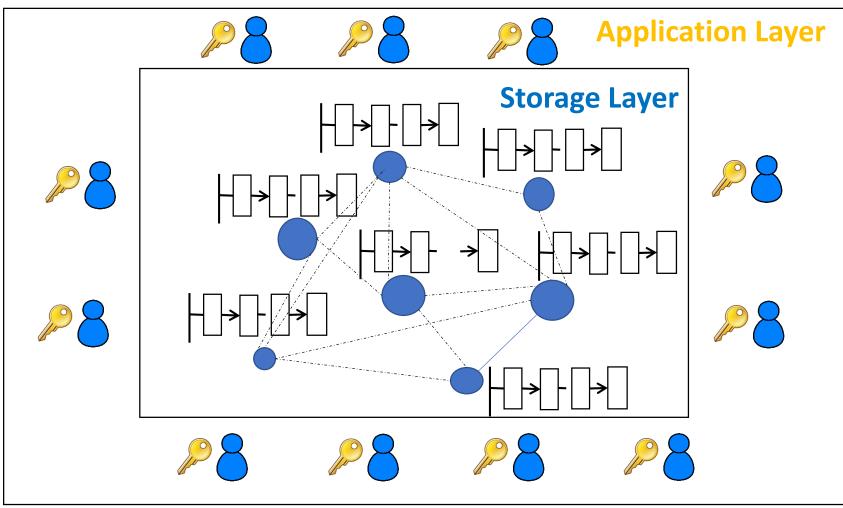
# Traditional Banking Systems

- From Database and Distributed Computing Perspective
- Identities and Signatures
  - You are your signature: IDENTITY
    Private and Public Digital signatures
- Ledger
  - The balance of each identity (saved in a DB)
  - → Blockchain (basically a linked list!)
- Transactions
  - Move money from one identity to another
  - Concurrency control to serialize transactions → Mining and Proof of Work
  - Typically backed by a transactions log
    - Log is persistent (disk) → Replication to the whole world
    - Log is immutable and tamper-free (end-users trust this) → HashPointers

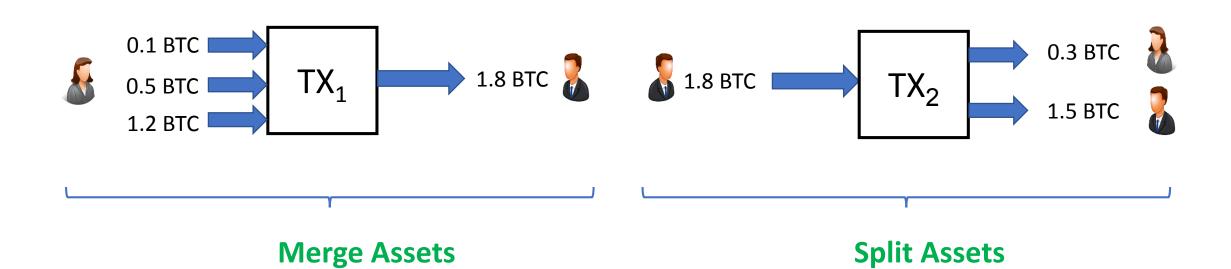
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# Blockchain Architecture

- The ledger is fully replicated to all network nodes
  - A Block is a set of transactions submitted by the clients.

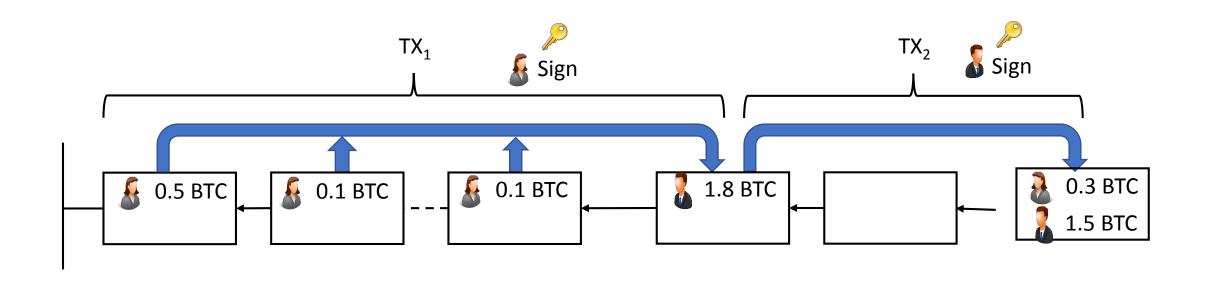


#### **Transaction Model**



Assuming no imposed transaction fees!

#### **Transaction Model**



#### The Ledger: Some Technical Details

#### • How is the ledger tamper-free?

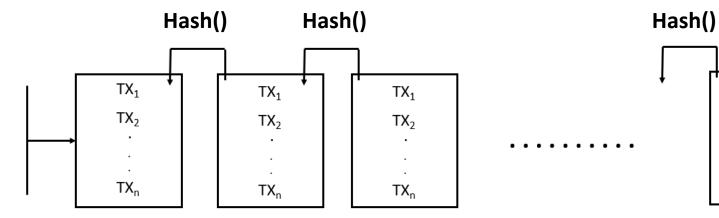
Blocks are connected through **hash-pointers** 

- Each block contains the hash of the previous block header
- Tampering with the content of any block can easily be detected

 $TX_1$ 

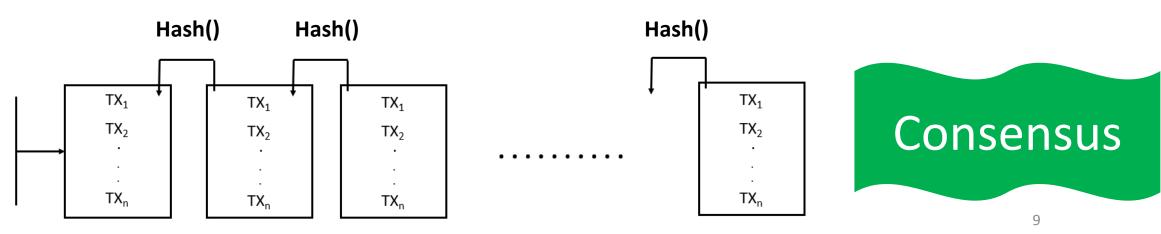
 $TX_2$ 

TX<sub>n</sub>



# Making Progress

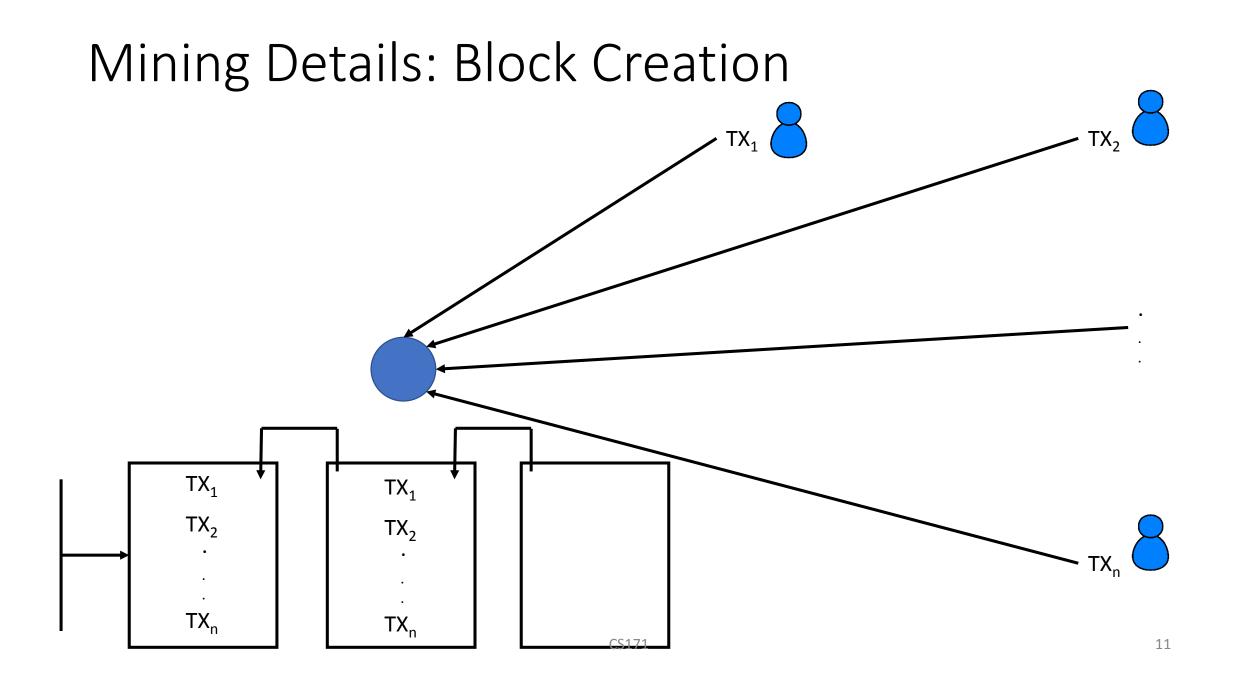
- To make progress:
  - Network nodes validate new transactions to make sure that:
    - Transactions on the new block do not conflict with each other
    - Transactions on the new block do not conflict with previous blocks transactions
  - Network nodes need to agree on the next block to be added to the blockchain
- New assets are generated and registered through mining.
  - Reward transaction in every mined block



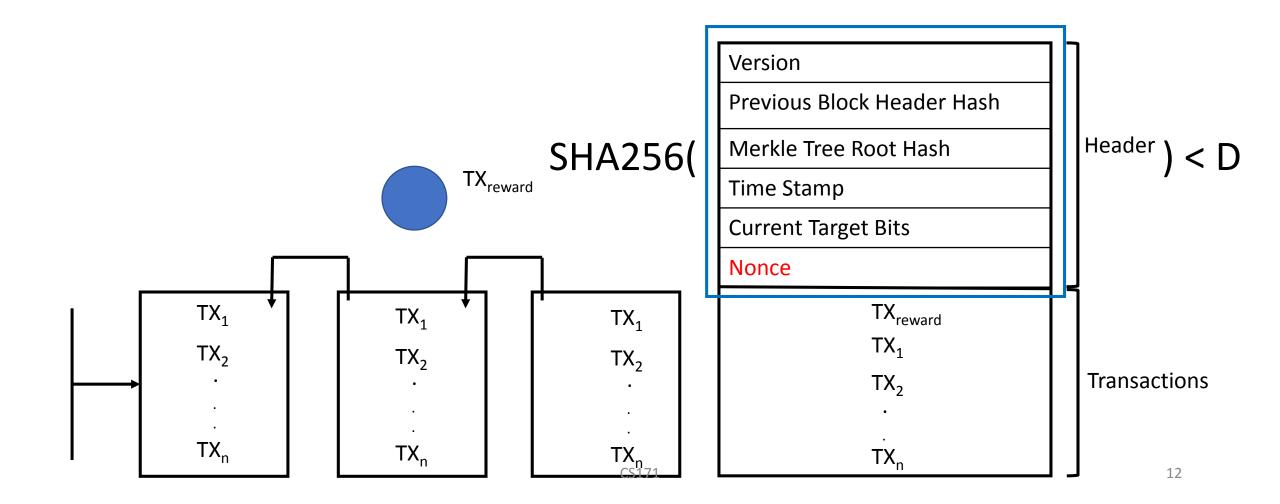
# **Consensus Protocols**

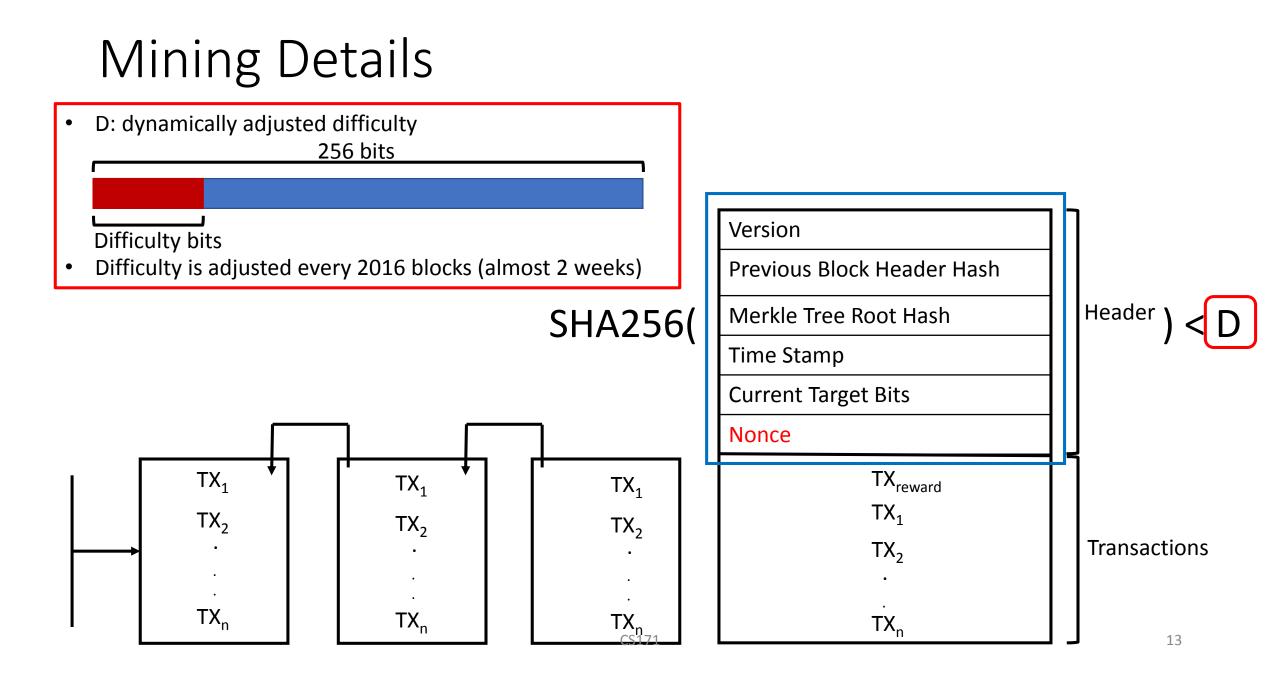
All participants should be known a priori

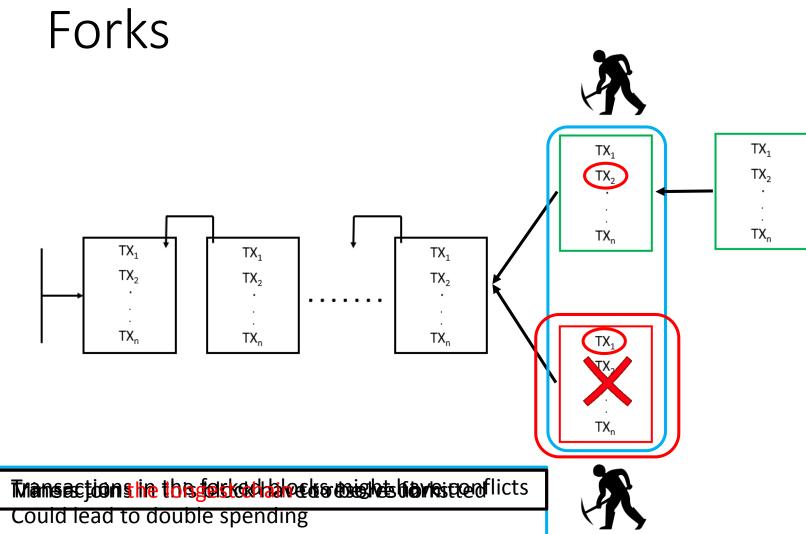
- Permissioned vs Permissionless settings
- Permissionless Blockchains:
  - Network nodes freely join or leave at anytime
  - Nakamoto's Consensus: Proof of Work (PoW)
  - Ethereum's Consensus: Proof of Stake (PoS)
- : Permissioned Blockchains
  - Paxos (Crash failures only)
  - Byzantine Fault-tolerance (malicious failures)



#### Mining Details: Block Contents







• Forks have to be eliminated

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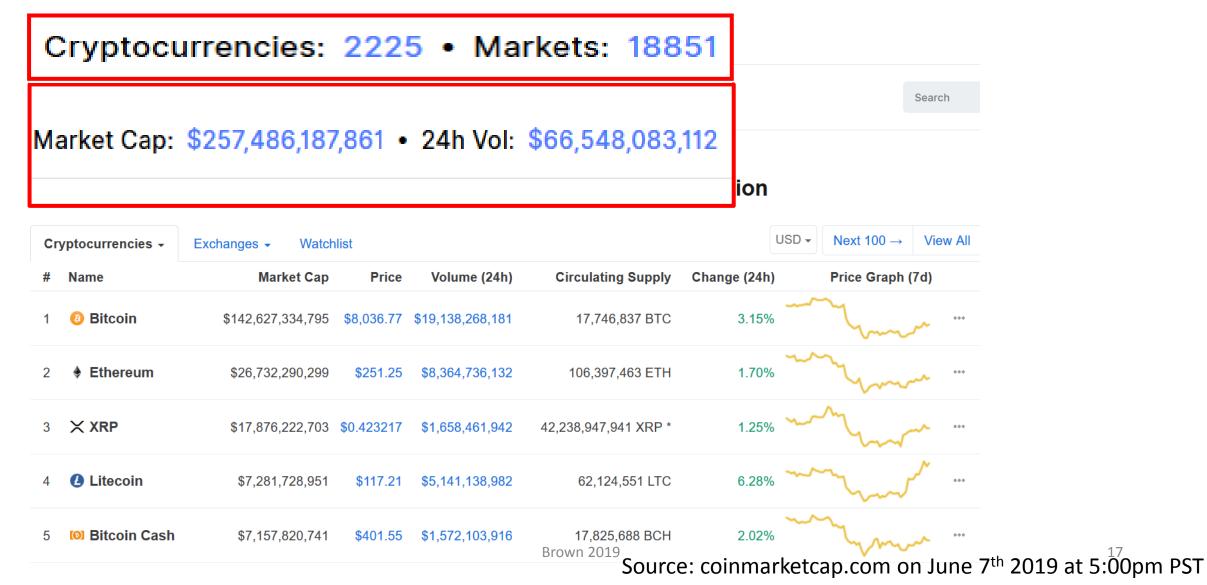
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#### Some Limitations of Bitcoin

- High transaction-confirmation latency
- Probabilistic consistency guarantees
- Very low TPS (Transactions per second) average of **3 to 7 TPS**
- Transparency leads to lack of privacy
- Energy consumption due to PoW.

# Atomic Commitment Across Blockchains

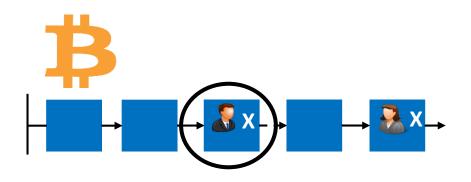
#### The Landscape

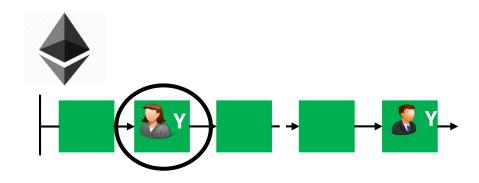


# The Landscape

- Thousands of Blockchains
- Tens of thousands of markets
- Exchanges to trade tokens for USD
- Direct token transactions in one blockchain
- Direct token transactions across blockchains, how?
- Cross-chain transactions

#### Cross-ChainTransaction Example





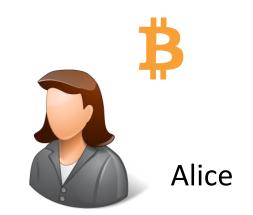
# Atomic Cross-Chain Commitment Protocol



Swap of Ownership

• Alice wants to trade Bitcoin for Ethereum with Bob



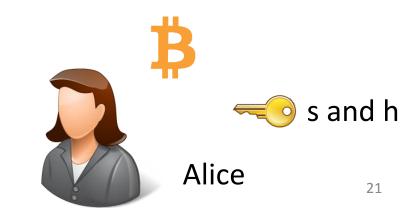


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• Alice wants to trade Bitcoin for Ethereum with Bob



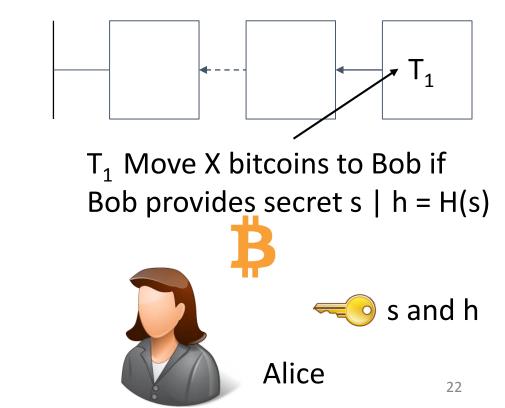
- Create a secret s —
- Calculate its hash h = H(s)



Brown 2019

• Alice wants to trade X Bitcoin for Y Ethereum with Bob

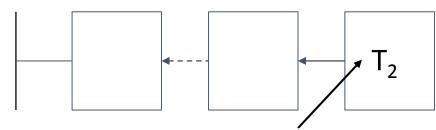
Bitcoin blockchain





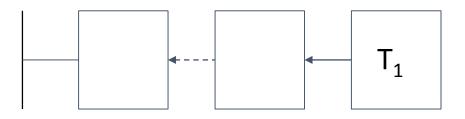
• Now, h is announced in Bitcoin blockchain and made public

Ethereum blockchain



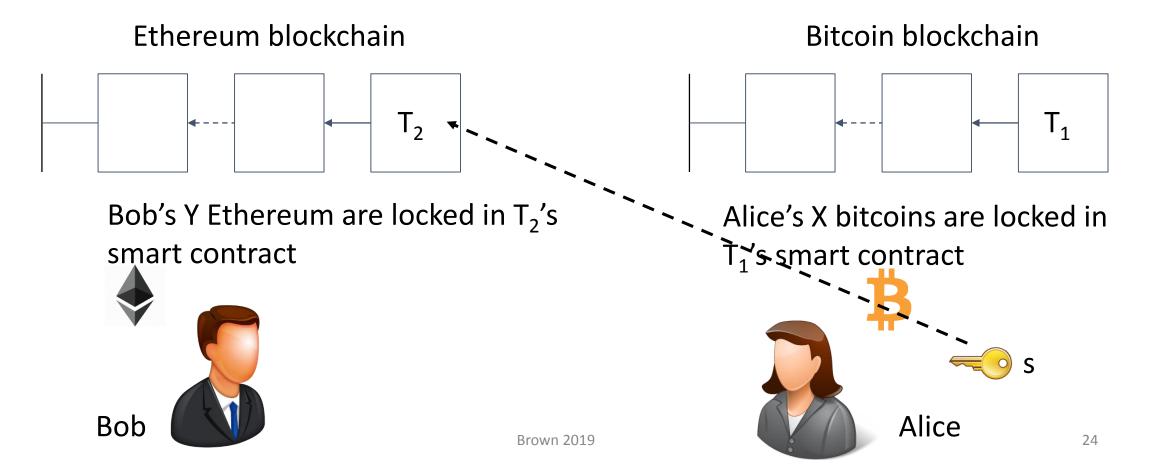
Bob

 $T_2$  Move Y Ethereum to Alice if Alice provides secret s | h = H(s) Bitcoin blockchain



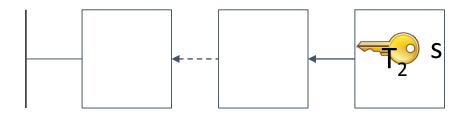
Alice's X bitcoins are locked in T<sub>1</sub>'s smart contract

• Now, for Alice to execute T<sub>2</sub> and redeem Y Ethereum, she reveals s



• Revealing s, executes T<sub>2</sub>. Now s is public in Ethereum's blockchain

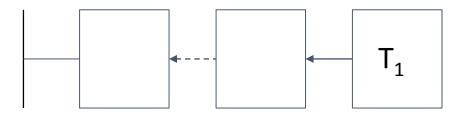
Ethereum blockchain



Bob's Y Ethereum are locked in  $T_2$ 's smart contract

Bob

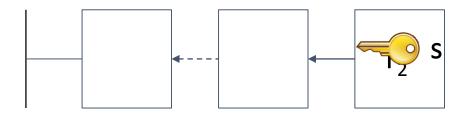
Bitcoin blockchain



Alice's X bitcoins are locked in T<sub>1</sub>'s smart contract Alice 25

• Now, Bob uses s to execute  $T_1$  and redeem his Bitcoins

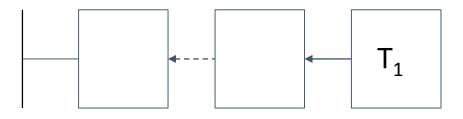
Ethereum blockchain



Bob's Y Ethereum are locked in  $T_2$ 's

smart contract Bob

Bitcoin blockchain



Alice's X bitcoins are locked in T<sub>1</sub>'s smart contract Alice

# Atomic Swap Example: What can go wrong?

- Alice locks her X Bitcoins in Bitcoin's blockchain through T<sub>1</sub>
- Bob sees T<sub>1</sub> but refuses to insert T<sub>2</sub>
- Now, Alice's Bitcoins are locked for good
  - A conforming party (Alice) ends up worse off because Bob doesn't follow the protocol
- Prevention
  - Use timelocks to expire a contract
  - Specify that an expired contract is refunded to the creator of this contract

### Atomic Swap Example: Timelocks

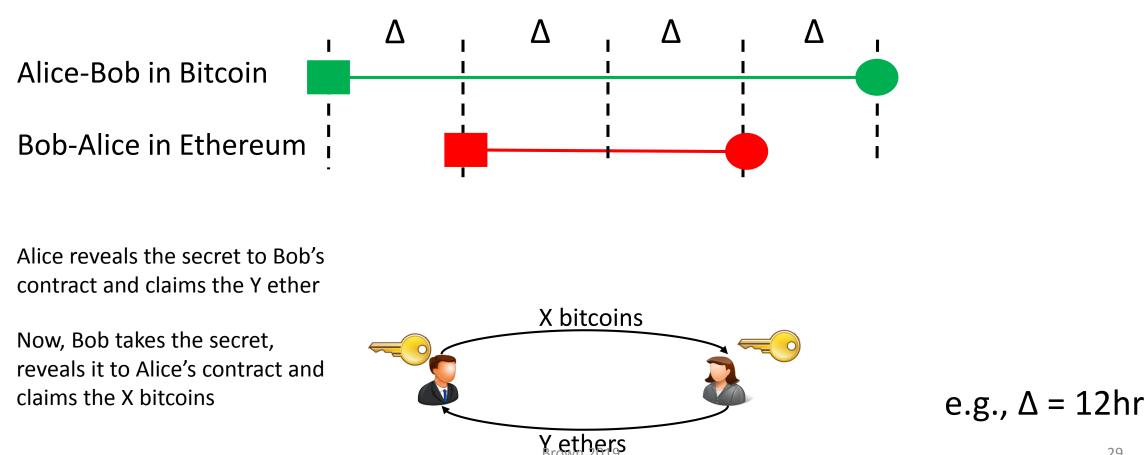
How to determine the time period of a timelock?

 $T_4$ : Refund  $T_2$  to Bob if Alice does not execute  $T_2$  before **24** hours

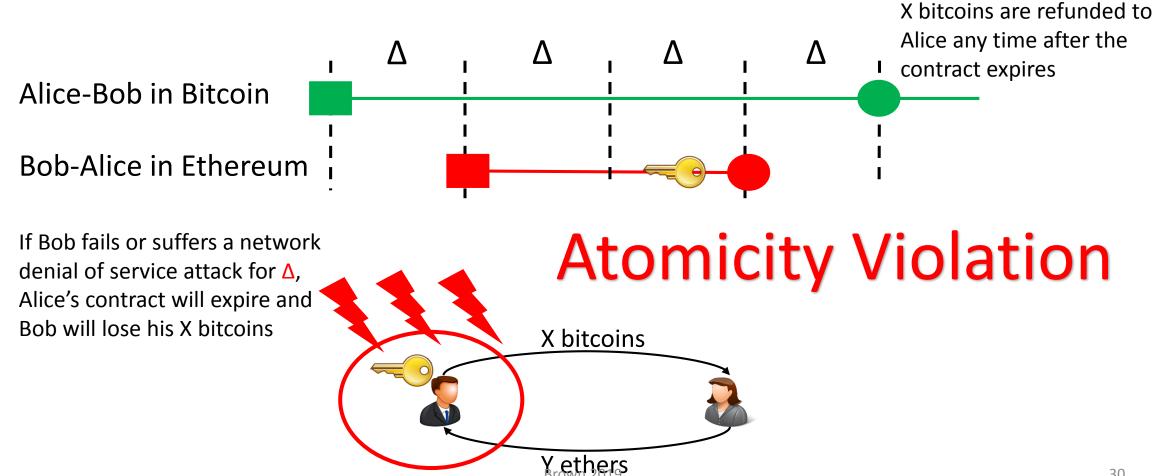
 $T_2$ : Move Y Ethereum to Alice if Alice provides secret s | h = H(s)



T<sub>3</sub>: Refund T<sub>1</sub> to Alice if Bob does not execute T<sub>1</sub> before 48 hours T<sub>1</sub>: Move X bitcoins to Bob if Bob provides secret s | h = H(s)



# What can go wrong?



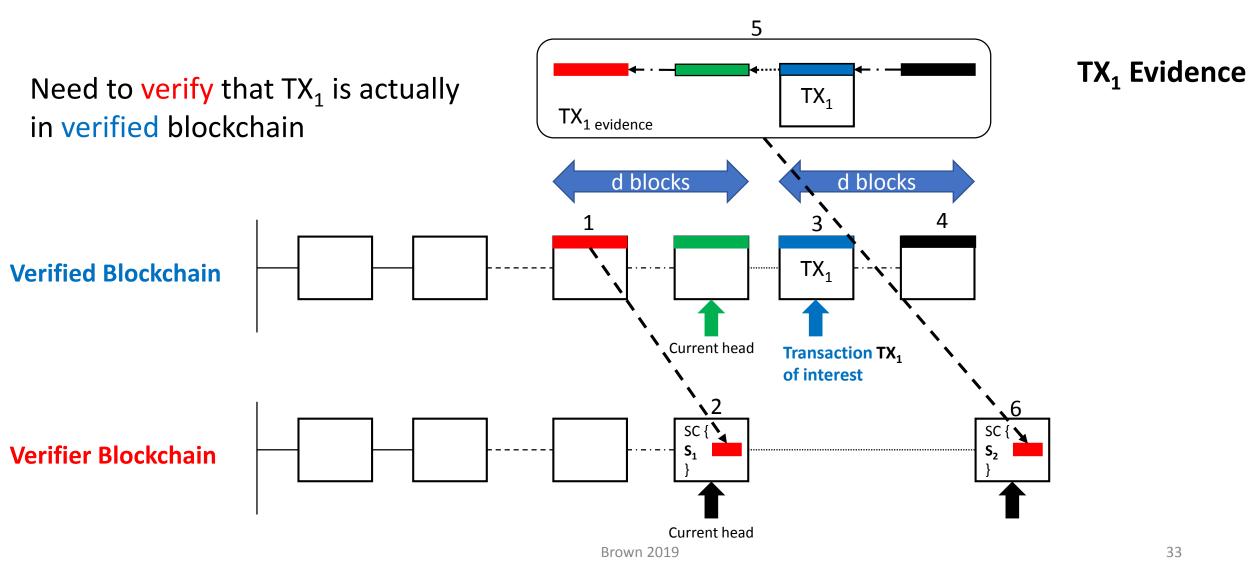
#### Atomicity Violation

- Using timelocks leads to Atomicity violation
- Our Atomicity-based Approach:
  - The decision of both transactions should be made atomic
    - Once the decision is taken, both transactions either commit or abort
  - A transaction cannot commit unless a commit decision is reached
  - A transaction cannot abort unless an abort decision is reached

#### Building block: Cross-Chain Verification

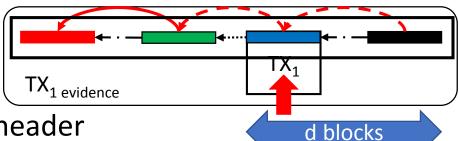
- How can miners of one blockchain:
  - Verify a transaction in another blockchain?
  - Without maintaining a copy of this other blockchain.

#### Building block: Cross-Chain Verification



#### Building block: Cross-Chain Verification

- Verification process:
  - Each header includes the hash of the previous header
  - The proof of work of each header is correct
  - TX<sub>1</sub> is correct
  - TX<sub>1</sub> is buried under d blocks
- The cost of generating evidence:
  - Choose d to make this cost > the value transacted in TX<sub>1</sub>
  - If true, a malicious user has no incentive to create a fake evidence

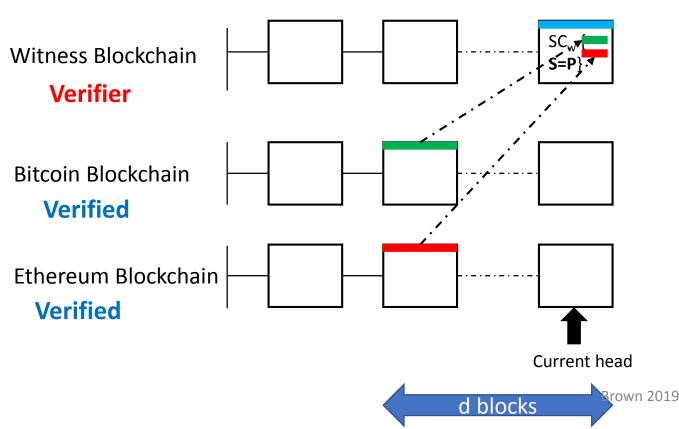


#### Atomic Commitment Across Blockchains

- Use another blockchain to witness the Atomic Swap
- The witness blockchain decides the commit or the abort of a swap
- Once a decision is made:
  - All sub-transactions in the swap must follow the decision
  - Achieves atomicity, either all committed or all aborted
- Cross chain verification is leveraged twice
  - Miners of the witness network verify the publishing of contracts in asset blockchains
  - Miners of assets' blockchains verify the decision made in the witness network

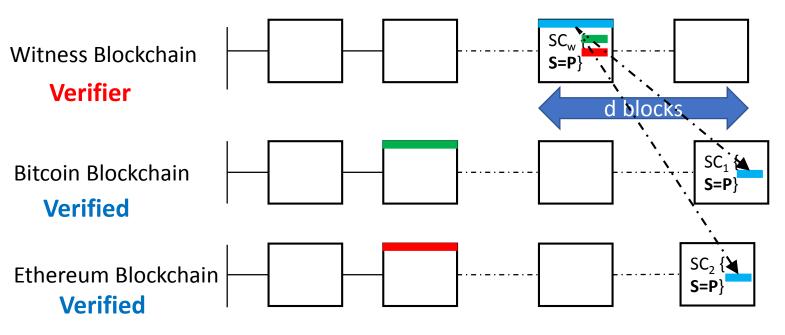
#### Protocol Sketch

- Deploy a contract SC<sub>w</sub> in the witness network with state Published (P)
- $SC_w$  has a header of a block at depth d of all blockchains in the swap



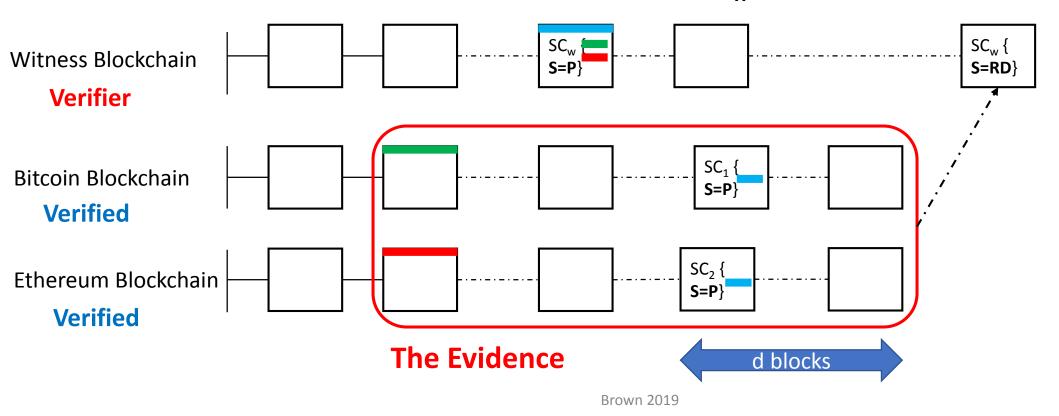
#### Protocol Sketch Cont'd

- Participants deploy their contracts in the corresponding blockchains
- Participants add the header of SC<sub>w</sub> to their contracts



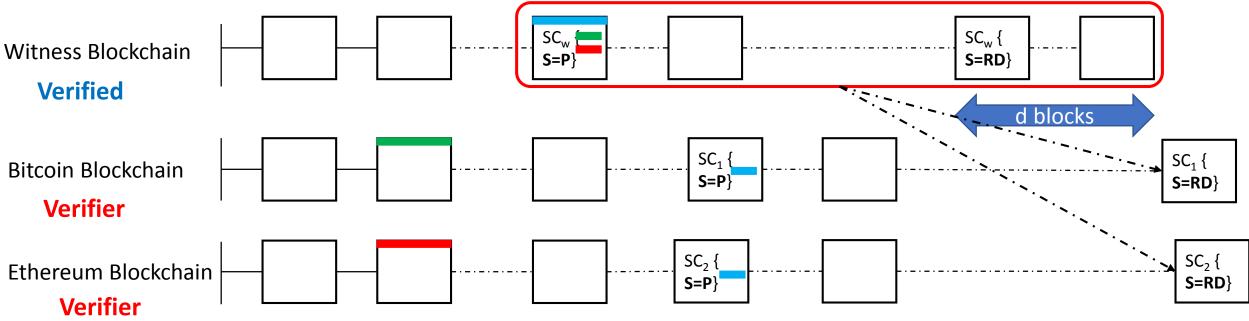
### Protocol Sketch Cont'd

- Participants submit evidence of publishing the smart contracts in Assets Blockchains
- If all contracts are published and correct, SC<sub>w</sub>'s state is altered to redeem (RD)



### Protocol Sketch Cont'd

- Participants submit evidence of Redeem State (RD) from the Witness Blockchain to the Assets Blockchains.
- After evidence verification, participants redeem their assets from the Assets Blockchains.



#### Atomic Commitment Across Blockchains

- SC<sub>w</sub>'s state determines the commit (RD) or the abort (RF) decision
- Once SC<sub>w</sub>'s state is altered and the block is buried under d blocks:
  - All sub-transactions must follow this decision
  - None of the sub-transactions can decide on a different decision
- Even if a participant fails or faces a network denial of service:
  - When the participant recovers, the evidence of the decision still exists
  - This evidence can be used to redeem or refund the contracts
- The only way to violate atomicity is to fork the witness blockchain
- Economic incentives prevent this attack
- Any protocol is prone to fork attacks

# Parting Thoughts

• Building global-scale blockchains is a collective effort.

